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Kekulé Count in Toroidal Hexagonal Carbon Cages

Peter E. John

Technical University Ilmenau, Institute of Mathematics, PSF 100565, D-98684 Ilmenau, Germany

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After the fullerenes have been found, it is a natural question to ask whether there are torus-shaped "graphitoid" carbon molecules which may be called $toroidal\ graphitoids$. Note that the torus is the only closed surface S that can carry graphs G such that all vetices of G have degree 3 and all faces of the embedding of G in S are hexagons (Figure 1). In what follows, such (hypothetical) molecules (see Ref. 1a) and their molecular graphs will be referred to as "torenes".

Note that the first paper about this topic was given by M. Randić, Y. Tsukano, and H. Hosoya. ^{1b}

In this paper, an algorithm is given that enables the number of Kekulé structures of a torene to be calculated in polynomial time (the complexity problem will be discussed elsewhere).



Figure 1

TOROIDAL HEXAGONAL TESSELLATIONS (TORENES)

Let **T** denote the class of all torenes. For $T \in \mathbf{T}$ let h = h(T), v = v(T) and e = e(T) denote the numbers of hexagons, vetices and edges of T, respectively. Clearly, v = 2h and e = 3h. The graph of $T \in \mathbf{T}$ can be drawn in the

plane (equipped with the regular hexagonal lattice L) using the representation of the torus T by a parallelogram P with the usual boundary identification (see Figure 2, e.g., parallelogram $A = A_1$, A'_1 , B''_1 , B'''_1 – note that the points π_1 , π'_1 , π_2 and π'_2 represent the same point of T). Let d_i (i = 1,2,3) be the edge directions of L (Figure 2).

Select $P=P_i$ such that its sides s_i , s'_{i-1} are perpendicular to d_i , d_{i-1} , respectively, (the subscript i-1 is to be reduced to the smallest positive integer modulo 3). Let p_i and q_i denote the number of hexagons met by s_i , and the number of layers of hexagons parallel to s_i that are covered by P, respectively, and let t_i (the torsion) denote the number of edges of L intersected by s_i between A and B_i . In this way, for every $T \in \mathbf{T}$, three (not necessarily incongruent) representations T_1 , T_2 , T_3 of its graph are obtained.

W.l.o.g. we may assume $p_1 \leq p_2 \leq p_3$. Let $p = p_1$, $q = q_1$, $t = t_1$, $s = s_1$, $T^* = T_1$: then T^* and p, q, t may be considered the canonical representation, and the canonical parameters, of the graph $T \in \mathbf{T}$, respectively. We shall briefly write $T^* = (p, q, t)$. Note that, for fixed p, q, t, all torenes T with parameters p, q, t where $t' \equiv t \pmod{p}$ have the same canonical representation $T^* = (p, q, t)$. For our example given in Figure 2, clearly p = 5, q = 6 and t = 1. The canonical parameters can easily be found using, e.g., the method described in Ref. 2.

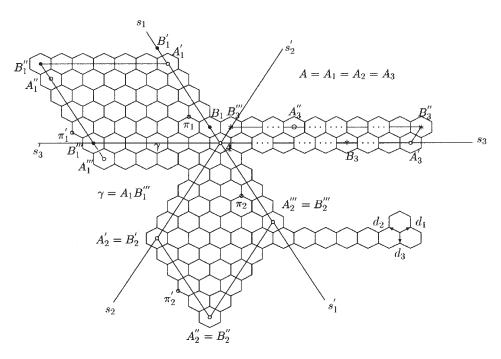


Figure 2

DEFINITIONS

Let G = (V, E) be a connected graph with vertex set V = V(G) and edge set E = E(G). For G, v = v(G) = |V| and e = e(G) = |E| denote the number of vertices and the number of edges of G respectively.

G is called *bipartite* iff its vertices can be coloured white and black such that every edge connects a white vertex with a black one.

A matching M of G is a set of pairwise disjoint edges of G. Matching M is called perfect iff M covers all vertices of G. Let $\mathbf{M} = \mathbf{M}(G)$ denote the set of all perfect matchings of G; set $m = m(G) = |\mathbf{M}|$. Let $M \in \mathbf{M}$. The edges of G that belong (do not belong) to M are referred to as the red (blue) edges. The number of perfect matchings of G that contain (do not contain) a given edge $k \in E$ is denoted by r(G, k) (b(G, k)): thus r(G, k) + b(G, k) = m(G) for every edge $k \in E$.

Note that, for $i \in \{1,23\}$, the set of all edges of direction d_i is a perfect matching of T.

PRELIMINARIES

Assume that T^* has been drawn such that s lies horizontally in the plane (Figure 3). Graph T^* being bipartite, its vertices can be coloured such that every vertical edge connects a black top vertex with a white bottom vertex. Let $\mathbf{P} = \mathbf{P}(T^*)$ denote the set of all vertical edges of T^* intersected by s; clearly, $|\mathbf{P}| = p$. Consider the 2^p subsets \mathbf{R} of \mathbf{P} where $|\mathbf{R}| = r \in \{0,1,...,p\}$. Set $\mathbf{B} = \mathbf{P} - \mathbf{R}$.

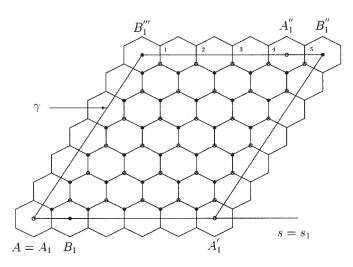


Figure 3

Let $\mathbf{M}(\mathbf{R})$ denote the set of perfect matchings of T that contain all edges of \mathbf{R} but no edge of \mathbf{B} (note that $\mathbf{M}(\mathbf{R})$ may be empty).

$$\mathbf{M}(T) = \bigcup_{\mathbf{R} \subseteq \mathbf{P}} \mathbf{M}(\mathbf{R}) . \tag{1}$$

Set

$$m(\mathbf{R}) = |\mathbf{M}(\mathbf{R})|. \tag{2}$$

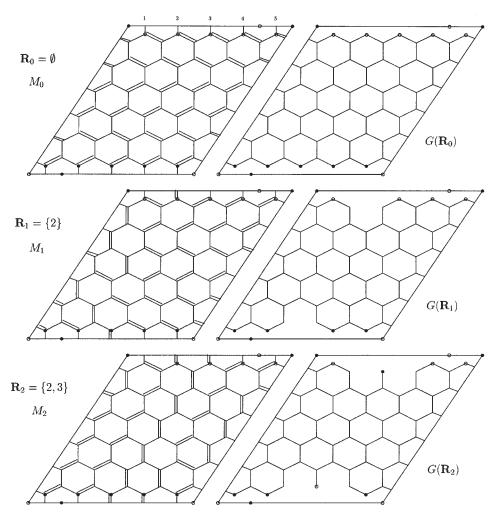


Figure 4.1. In the left-hand figures, one of the matchings in M(R) is depicted using double lines.

Graph $G(\mathbf{R})$ is obtained from T by omitting all edges of \mathbf{P} , all end vertices of the edges of \mathbf{R} , and all edges incident upon these vertices (Figure 4.1.). Set $\mathbf{G} = \{G(\mathbf{R}) | \mathbf{R} \subseteq \mathbf{P}\}$. Clearly, $m(G(\mathbf{R})) = m(\mathbf{R})$ thus by Eqs. (1) and (2),

$$m(T) = \sum_{\mathbf{R} \subseteq \mathbf{P}} m(\mathbf{R}) = \sum_{r=0}^{p} \sum_{|\mathbf{R}|=r} m(\mathbf{R})$$
 (3)

Direct all edges of $G \in \mathbf{G}$ from top to bottom, thus turning G into a directed (bipartite) graph \vec{G} (Figure 4.2). Every \vec{G} is the embedding of a planar

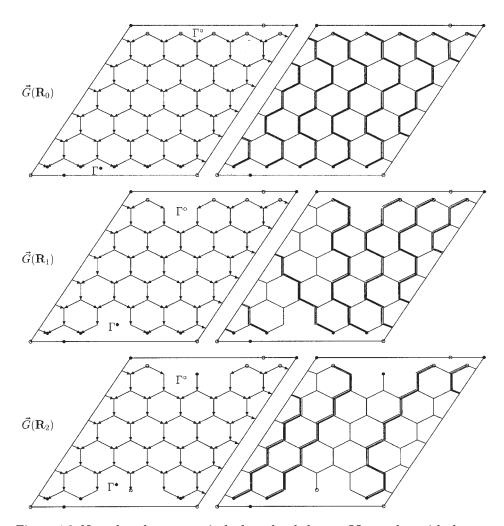


Figure 4.2. Note that the non-vertical edges that belong to $\mathbf{M_i}$ together with the vertical edges that do not belong to $\mathbf{M_i}$ (i = 0, 1, 2) form systems of disjoint monotone paths in $G(\mathbf{R}_i)$ connecting all peaks with all valleys of $G(\mathbf{R}_i)$; these path systems are depicted (see the heavy lines) in the figures on the right-hand side.

(not necessarily connected) directed graph in a cyclinder. Set $\vec{\mathbf{G}} = \{\vec{G} | G \in \mathbf{G}\}$. The white sources (black sinks) of $\vec{G} \in \vec{\mathbf{G}}$ are called its *peaks* (valleys). Note that a black source (white sink) is not a peak (valley). As in the case of hexagonal systems in the plane, $\vec{G} \in \vec{\mathbf{G}}$.

- (i) the number of peaks of a graph $\vec{G} = \vec{G}(\mathbf{R})$ is equal to the number of its valleys, namely, p r;
- (ii) for each $\mathbf{R} \subseteq \mathbf{P}$ there is a (1,1)-correspondence between the set $\mathbf{M}(\mathbf{R})$ and the set of systems of disjoint monotone paths connecting the peaks of $\vec{G}(\mathbf{R})$ with its valleys such that in every such path system each peak and each valley are the end vertices of exactly one (directed) path (Figure 4.2). Every $\vec{G} \in \vec{\mathbf{G}}$ is the embedding of a planar (not necessarily connected) directed graph in a cylinder.

Some immediate observations.

- (a) $\vec{G}(\mathbf{R})$ is the empty graph (without edges and vertices) if and only if $\mathbf{R} = \mathbf{P}$ and q = 1. By convention, the number of monotone path systems of the empty graph is assumed to be equal to 1.
- (b) If $\vec{G} = \vec{G}(\mathbf{R})$ is non-empty and connected, then it has two particular (in general, non-hexagonal) faces which result from the edge-and-vertex deleting process described above, namely a top face $\Gamma^{\circ} = \Gamma^{\circ}(\vec{G})$ with all sources of \vec{G} on its boundary and a bottom face $\Gamma^{\bullet} = \Gamma^{\bullet}(\vec{G})$ with all sinks of \vec{G} on its boundary (see Figure 4.2.).
- (c) If \vec{G} is disconnected with components \vec{G}_1 , \vec{G}_2 ,..., \vec{G}_c , then, necessarily, $1 \le q \le 2$, and $m(G) = m(G_1) \cdot m(G_2) \cdot \ldots \cdot m(G_c)$. (Note that some component G_γ may have an odd number of vertices, implying m(G) = 0.)

KASTELEYN'S ORIENTATION

Graph orientations of the kind to be considered next were first used by P. W. Kasteleyn⁵ in his pioneering work on dimer coverings of the square lattice graph. Let \vec{G} be any directed plane graph – *i.e.*, an embedding of a directed planar graph in the plane – with an even number of vertices. Let $\omega = \omega(\vec{G})$ denote the orientation of \vec{G} , and let Γ be a (finite or infinite) face of G. Orientation ω is said to be *Kasteleyn with respect to* Γ iff, for each component C of the boundary of Γ , the number of arcs of C whose left bank belongs to Γ is odd; ω is *Kasteleyn* for G iff ω is Kasteleyn with respect to all faces (including the infinite face) of G.

It is not difficult to prove that a (finite) plane graph G has a Kasteleyn orientation if and only if each component of G has an even number of vertices. Thus, if G does not have a Kasteleyn orientation, then m(G) = 0. Let G be a plane graph equipped with a Kasteleyn orientation and let $\mathbf{A} = (a_{i,j})$ denote its Kasteleyn adjacency matrix defined by

1 if there is an are from vertex i to vertex j, $a_{ij} = \begin{cases} -1 & \text{if there is an arc from vertex } j \text{ to vertex } i, \\ 0 & \text{otherwise.} \end{cases}$

Kasteleyn's Theorem. 5,6

$$m(G) = \sqrt{\det \mathbf{A}} \ . \tag{4}$$

For $\mathbf{R} \subseteq \mathbf{P}$, besides $\omega_1 = \omega(\vec{G}(\mathbf{R}))$, consider the orientation ω_2 that is obtained from ω_1 by reversing the direction of all arcs intersected by the line segment $\gamma = AB_1^{""}$ of P (Figures 2, 3). Clearly, both orientations ω_1 , ω_2 are Kasteleyn with respect to every hexagon of every graph $G(\mathbf{R})$, $\mathbf{R} \subseteq \mathbf{P}$; in addition, it is easy to check that ω_{ρ} is Kasteleyn with respect to both Γ° and Γ^{\bullet} if and only if $p - r \equiv \rho$, mod 2. Define $\lambda = \lambda(\mathbf{R})$ by $\lambda \in \{1, 2\}$, $\lambda \equiv p - r$, mod 2. The above now yields

Lemma 1.

Orientation $\omega_{\lambda(\mathbf{R})}$ is Kasteleyn for $G(\mathbf{R})$.

For $\lambda = 1$, 2 and for every directed edge (arc) $\alpha \in E(\vec{G})$ define $\sigma_{\lambda}(\alpha)$ by

$$\sigma_1(a) = 1, \ \sigma_2(a) = \begin{cases} -1 & \text{if } a \text{ is intersected by } \gamma \\ 1 & \text{otherwise} \end{cases}$$

CALCULATING THE NUMBER OF KEKULÉ STRUCTURES

For $T \in \mathbf{T}$ we shall now determine the numbers $m(\mathbf{R})$ and m(T) (see equation (3)).

We start with $\mathbf{R}=\emptyset$ and set $\vec{G}(\emptyset)=\vec{G}_0$. Let $\left\{x_1,\,x_2,...,\,x_p\right\}$ and $\left\{y_{1-t},\,y_{2-t},...,\,y_{p-t}\right\}$ (labelled from left to right), be the set of peaks and the set of valleys of G_0 , respectively. Recall that t stands for the torsion, the subscripts 1-t, 2-t,..., p-t are to be reduced to the smallest positive residues modulo p.

Algorithm A (see Figure 5):

The entries are arranged as indicated in Figure 5a. For $\lambda = 1, 2$ assign to every vertex z of \vec{G}_0 a vector $\boldsymbol{w}_{\lambda}(z) = (w_{\lambda 1}(z), w_{\lambda 2}(z), ..., w_{\lambda p}(z))$ according to the following rules.

- (A.1)Start with the sources. For peak x_i (i = 1, 2,..., p), set $\boldsymbol{w}_{\lambda}(x_i) = (\delta_{i1}, \delta_{i2},..., \delta_{ip})$ where $\delta_{ii} = 1$ and $\delta_{ij} = 0 \text{ iff } i \neq j.$ For a black source y, set $\mathbf{w}_{\lambda}(y) = (0, 0, ..., 0) (\lambda = 1, 2)$.
- For every vertex z of \vec{G}_0 that is not a source, following the arrows of (A.2) G_0 calculate the vectors

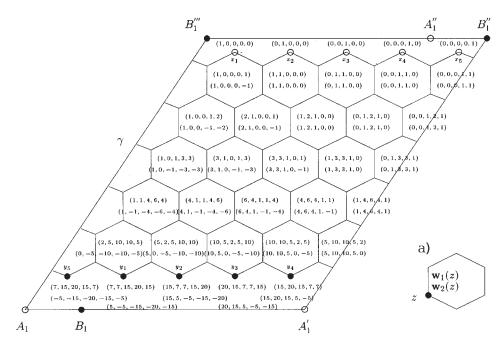


Figure 5. The entries are arranged as indicated in Figure 5a.

$$\mathbf{w}_{\lambda}(z) = \sum \sigma_{\lambda}(z', z) \cdot \mathbf{W}_{\lambda}(z'), (\lambda = 1, 2),$$

where the sum is taken over the one or two predecessors z' of z.

(A.3) For valley
$$y_j$$
, $j = 1, 2,..., p$, let $w_{\lambda,ij} = w_{\lambda i}(y_j)$ and $\boldsymbol{w}_{\lambda}(y_j) = (w_{\lambda 1j}, w_{\lambda 2j},..., w_{\lambda,pj})$. Form the $p \times p$ -matrices

$$\mathbf{W}_{\lambda}(t) = \mathbf{W}_{\lambda}(\vec{G}_{0})$$

$$= (\mathbf{w}_{\lambda}^{\mathrm{T}}(y_{1}), \ \mathbf{w}_{\lambda}^{\mathrm{T}}(y_{2}), ..., \ \mathbf{w}_{\lambda}^{\mathrm{T}}(y_{p}))^{\mathrm{T}}$$

$$= (\mathbf{w}_{\lambda, ij}),$$

$$i, j = 1, 2, \dots p \ (\lambda = 1, 2).$$

It is not difficult to see that the matrices $\mathbf{W}_1(0)$, $\mathbf{W}_2(0)$ resulting from the above algorithm can also be written as powers of the circulant $Z_+ = circ$ (1, 0,..., 0, 1) and the skew circulant $Z_- = scirc$ (1, 0,..., 0, -1), respectively, namely, $\mathbf{W}_1(0) = Z^q_+$, $\mathbf{W}_2(0) = Z^q_-$ (see Ref. 7). From $\mathbf{W}_{\lambda}(0)$ matrix $\mathbf{W}_{\lambda}(t)$ is obtained just by rotating (i.e., cyclically permuting) the rows of $\mathbf{W}_{\lambda}(0)$: the (t+1)st row of $\mathbf{W}_{\lambda}(0)$ multiplied by $(-1)^{\lambda-1}$ becomes the first row of $\mathbf{W}_{\lambda}(t)$, etc.

(A.4) For $\mathbf{R} \subseteq \mathbf{P}$, set $r = |\mathbf{R}|$ ($0 \le r \le p$). Clearly, the graph $G(\mathbf{R})$ is an induced subgraph of $G(\emptyset)$ with precisely p - r peaks and as many valleys. Let $\mathbf{W} = \mathbf{W}(\mathbf{R})$ denote the principal submatrix of $\mathbf{W}_{\lambda(\mathbf{R})} = \mathbf{W}_{\lambda(\mathbf{R})}(t)$ that corresponds to the complement of \mathbf{R} . Note that $\mathbf{W}(\mathbf{R})$ is a square matrix of order p - r, obtained from $\mathbf{W}_{\lambda(\mathbf{R})}$ by deleting all those rows and columns that correspond to the elements of \mathbf{R} . By Kasteleyn's theorem (formula (4), specified for bipartite graphs),

Lemma 2.

$$m(\mathbf{R}) = |\det \mathbf{W}(\mathbf{R})|. \qquad \Box (5)$$

The extreme cases

1. $\mathbf{R} = \mathbf{P}, p - r = 0.$

 $G(\mathbf{P})$ has no peak and no valley, $W(\mathbf{P})$ is the empty matrix. $G(\mathbf{P})$ has exactly one perfect matching which consists of all vertical edges (all edges perpendicular to s). Thus, in accordance with Eq. (5),

$$m(\mathbf{P}) = |det \mathbf{W}(\mathbf{P})| = 1$$
.

2. **R** = \emptyset , p - r = p.

In this simple case, we do not need Lemma 2: by immediate inspection we see that none of the vertical edges can belong to a perfect matching of $G(\emptyset)$ implying

$$m(\emptyset) = s^q$$
.

Equations (3) and (5) now allow m = m(T) to be calculated:

$$m(T) = \sum_{\mathbf{R} \subset \mathbf{P}} m(\mathbf{R}) = \sum_{\mathbf{R} \subset \mathbf{P}} |det \ \mathbf{W}(\mathbf{R})| \ . \tag{6}$$

Let $\mathbf{P} = \{1, 2, ..., p\}$ be the set of edges intersected by s where edge i is incident upon peak v_i (see Figures 3).

Call two sets \mathbf{R} , $\mathbf{R'} \subseteq \mathbf{P}$ cyclically equivalent (briefly: c-equivalent) iff $\mathbf{R'}$ is obtained from \mathbf{R} by shifting the subscripts defining \mathbf{R} modulo p. Clearly, for all members \mathbf{R} of a c-equivalence class the graphs $G(\mathbf{R})$ are isomorphic, thus it suffices to calculate $m(\mathbf{R}) = |det \mathbf{W}(\mathbf{R})|$ for only one representative \mathbf{R} of each c-equivalence class C and multiply by $\mu = \mu(\mathbf{R}) = |C|$.

Note that the above algorithm is applicable for any (not necessarily canonic) representation (p_i, q_i, t_i) (i = 1, 2, 3) of T.

For our example, the calculations and results are summarized in Table I, which is self-explanatory.

To make the situation more transparent, m(T) is calculated a second time in Table II, where another (non-canonical) representation of T-namely with $p_2 = 6$, $q_2 = 5$, $t_2 = 0$ (see Table II) — is used.

TABLE I (see Figure 3)

\mathbf{R}	r	μ	λ	W	m	$\mu \cdot m$
Ø	0	1	1	$oldsymbol{W}_1$	$2^6 = 64$	64
{1}	1	5	2	$ \begin{pmatrix} 5 & -5 & -15 & -20 \\ 15 & 5 & -5 & -15 \\ 20 & 15 & 5 & -5 \\ -15 & -20 & -15 & -5 \end{pmatrix} $	3125	15625
{1,2}	2	5	1	$\begin{pmatrix} 7 & 7 & 15 \\ 15 & 7 & 7 \\ 20 & 15 & 7 \end{pmatrix}$	1128	5640
{1,3}		5	1	$\begin{pmatrix} 7 & 15 & 20 \\ 20 & 7 & 7 \\ 15 & 15 & 7 \end{pmatrix}$	2983	14915
{1,2,3}	3	5	2	$\begin{pmatrix} 5 & -5 \\ -15 & -5 \end{pmatrix}$	100	500
{1,2,4}		5	2	$\begin{pmatrix} 5 & -15 \\ 20 & -5 \end{pmatrix}$	325	1625
{1,2,3,4}	4	5	1	(7)	7	35
$\{1,2,3,4,5\}$	5	1	2	Ø	1	1
Σ		$2^5 = 32$	/	/	/	m(T) = 38405

TABLE II

\mathbf{R}	r	μ	λ	$oldsymbol{W}$	m	$\mu \cdot m$
Ø	0	1	2	$oldsymbol{W}_2$	$2^5 = 32$	32
{1}	1	6	1	$\begin{pmatrix} 1 & 1 & 5 & 10 & 10 \\ 5 & 1 & 1 & 5 & 10 \\ 10 & 5 & 1 & 1 & 5 \\ 10 & 10 & 5 & 1 & 1 \\ 5 & 10 & 10 & 5 & 1 \end{pmatrix}$	1296	7776
{1,2}	2	6	2	$\begin{pmatrix} 1 & -1 & -5 & -10 \\ 5 & 1 & -1 & -5 \\ 10 & 5 & 1 & -1 \\ 10 & 10 & 5 & 1 \end{pmatrix}$	621	3726

R	r	μ	λ	W	m	$\mu \cdot m$
Ø	0	1	2	W_2	$2^5 = 32$	32
{1,3}		6	2	$\begin{pmatrix} 1 & -5 & -10 & -10 \\ 10 & 1 & -1 & -5 \\ 10 & 5 & 1 & -1 \\ 5 & 10 & 5 & 1 \end{pmatrix}$	1896	11376
{1,4}		3	2	$\begin{pmatrix} 1 & -1 & -10 & -10 \\ 5 & 1 & -5 & -10 \\ 10 & 10 & 1 & -1 \\ 5 & 10 & -5 & 1 \end{pmatrix}$	2561	7683
{1,2,3}	3	6	1	$\begin{pmatrix} 1 & 1 & 5 \\ 5 & 1 & 1 \\ 10 & 5 & 1 \end{pmatrix}$	76	456
{1,2,4}		6	1	$\begin{pmatrix} 1 & 5 & 10 \\ 10 & 1 & 1 \\ 10 & 5 & 1 \end{pmatrix}$	396	2376
$\big\{1,\!2,\!5\big\}$		6	1	$\begin{pmatrix} 1 & 1 & 10 \\ 5 & 1 & 5 \\ 10 & 10 & 1 \end{pmatrix}$	396	2376
{1,2,5}		2	1	$\begin{pmatrix} 1 & 5 & 10 \\ 10 & 1 & 5 \\ 5 & 10 & 1 \end{pmatrix}$	976	1952
{1,2,3,4}	4	6		$\begin{pmatrix} 1 & -1 \\ 5 & 1 \end{pmatrix}$	6	36
{1,2,3,5}		6	22	$\begin{pmatrix} 1 & -5 \\ 10 & 1 \end{pmatrix}$	51	306
{1,2,4,5}		3	2	$\begin{pmatrix} 1 & -10 \\ 10 & 1 \end{pmatrix}$	101	303
{1,2,3,4,5}	5	6	1	(1)	1	6
{1,2,3,4,5,6}	6	1	2	Ø	1	1
\sum		$2^6 = 64$	/	/	/	m(T) = 38405

For T = (p, q, t) write m(T) = m (p, q, t).

In order to show that m(T) depends on the torsion t of T, we calculated m(5, 6, t) for t = 0, 1, 2, 3, 4 and obtained

m(5, 6, 0) = m(5, 6, 1) = 38405, m(5, 6, 2) = m(5, 6, 4) = 38405,m(5, 6, 3) = 39440 (see also Ref. 8).

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SAŽETAK

Broj Kekuléovih struktura u torusnim heksagonalnim ugljikovim kavezima

Peter E. John

Nakon otkrića fullerena, prirodno se postavilo pitanje postojanja čisto ugljikovih heksagonalnih kaveza u obliku torusa koje bismo mogli nazvati »torusnim grafitom«. Istaknimo da je torus jedina zatvorena površina S na koju su mogu smjestiti trovalentni grafovi G čija su sva lica heksagoni (Ref. 1a, Slika 1). Predloženo je da se takove molekule i pripadni im molekularni grafovi zovu torenima.

Napomenimo da su se ovom temom prvi puta bavili M. Randić, Y. Tsukano i H. Hosoya. $^{\rm 1b}$

Ovdje je prikazan algoritam koji omogućava izračunavanje broja Kekuléovih struktura torena u polinomskom vremenu.



Slika 1