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A PASSENGER FLOW ROUTING MODEL FOR HIGH-SPEED RAILWAY NETWORK IN DIFFERENT TRANSPORTATION ORGANIZATION MODES

ABSTRACT

Reasonable selection of passenger flow routes considering different transportation organization modes can meet the demands of adapting to large-scale high-speed railway networks and improving network efficiency. Passenger flow routing models are developed to find and optimize a set of passenger flow routes for a high-speed railway network considering different transportation organization modes. In this paper, we presented a new approach minimizing the operating costs, including traveling cost, cost of travel time differences between different lines, and penalties for the inter-line. The network was reconstructed to solve the directed graph with four nodes (node-in-up, node-in-down, nodes-out-up, and nodes-out-down) indicating one station. To tackle our problem, we presented an integer non-linear programming model, and direct passenger demand was guaranteed owing to volume constraints. Binary variables were introduced to simplify the model, and the algorithm process was optimized. We suggested a global optimal algorithm by Lingo 11.0. Finally, the model was applied to a sub-network of the Northeast China railway system. Passenger flow routes were optimized and the transportation organization mode was discussed based on passenger volume, traveling distance, and infrastructure.

KEY WORDS

high-speed railway; train planning; transportation organization mode; passenger flow routing; integer non-linear programming model;

1. INTRODUCTION

High-speed railway (HSR) experiences a period of rapid development in China. The length of high-speed railways in service has reached 25,000 kilometers, ranked number one in the world in 2017, and it will increase to over 30,000 kilometers in 2020. High-speed railway has provided more convenient services for passengers. 30% of HSR lines operate at a maximum speed of 300 km/h in China, depending on infrastructure conditions.

1.1 Features analysis of a large-scale HSR network

There are three typical features of a large-scale HSR network. The features will be elaborated by the example of HSR in China.

Larger scale network and longer operating mileage

Compared with Germany, France, Dutch, and Japan, China's high-speed railway network obviously has larger scale and longer operating mileage. China has the longest mileage and the largest scale high-speed railway network globally. The length of HSR in service has reached 25,000 kilometers, which is more than 60 percent of total HSR in the world.

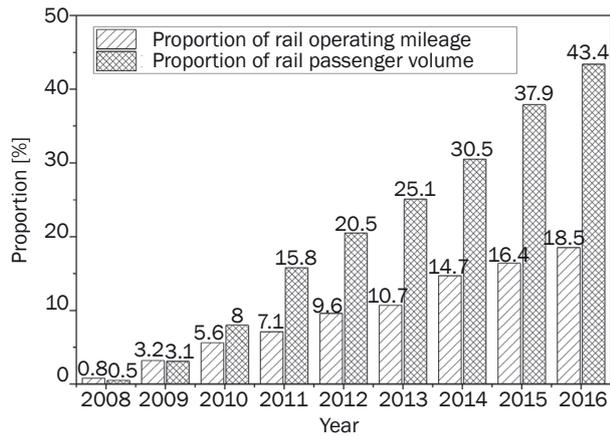


Figure 1 – High-speed railway change accounts for the proportion of railway passenger traffic volume and length of railways in operation

The HSR operating mileage and passenger volume are increasing year by year. High-speed railway takes 43.4% of the railway passenger volume with an operating mileage of 18.5% in China (Figure 1).

At the same time, there are HSR lines with a mileage of 1,000 km and above. For example, the length of the Beijing-Guangzhou high-speed railway is 2,298 kilometers, and it takes about 10 hours with a high-speed train from Beijing to Guangzhou.

Lower proportion of direct transportation and unbalanced passenger volume per direction

Since China has a vast territory and varied terrain, passengers have created a pressing demand for direct transport. Statistics shows that direct passenger train services were provided to only 10% OD pairs in 2016.

On the other hand, HSR provides lower train frequency for most OD pairs. 79.15% OD pairs are directly serviced with up to 5 trains, while 92.13% are directly serviced with up to 10 trains (Figure 2).

Unbalanced passenger volume by direction is widespread in large scale HSR networks. Figure 3 shows passenger volumes of several sections in the Baoji-Lanzhou HSR. There is a big passenger volume difference for some sections per directions.

Complex passenger transportation in hubs

There were 687 high-speed railway stations in China as of February 2017. Some cities have two or more high-speed railway stations. There is little difference in train classes between these stations. Connections through several stations in a single city add complexity to train operations. High-speed railway hubs have been formed with different scale and different features.

For example, there are seven main railway stations in Beijing: Beijing Railway Station, Beijing West Railway Station, Beijing South Railway Station, Beijing North

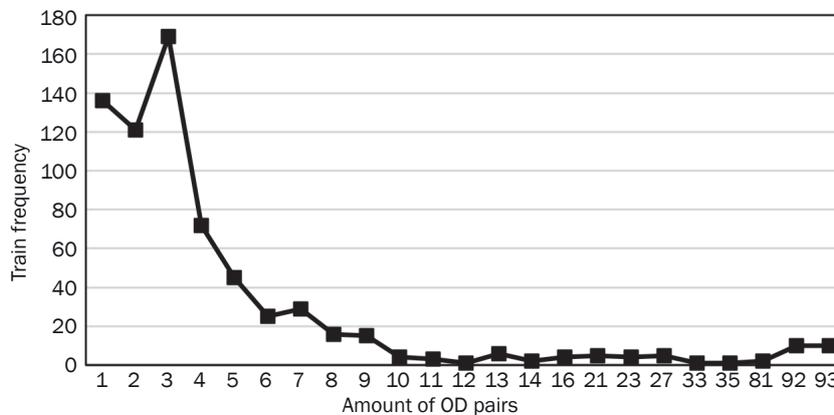


Figure 2 – The relationship of amount of OD pairs and train frequency

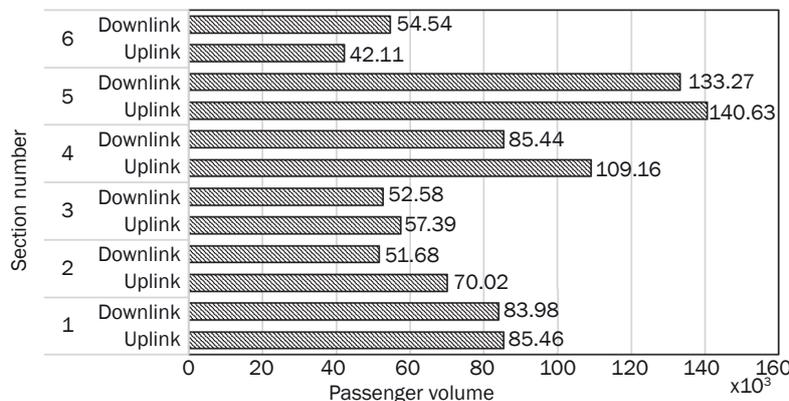


Figure 3 – Passenger volume of several sections per directions

Railway Station, Fengtai Railway Station, Xinghuo Railway Station, and Tongzhou Railway Station. All of the trains at the Beijing South Railway Station are EMUs, while Beijing West Railway Station has 60% EMUs. Other stations only have a few EMUs. Most of the stations cannot be connected to each other by railway.

More and more HSR stations will be constructed and operated in other cities with the development of the HSR network, such as Huizhou, Guangdong, which will have 9 stations.

1.2 Transportation organization mode analysis of a large-scale HSR network

These features have a big influence on the transportation organization mode of HSR, and the content of the transportation organization mode refers to direct or transfer, periodic or non-periodic, and intra-line or inter-line.

Definition of intra-line and inter-line. Intra-line: passengers' origin and destination belong to same line, and these passengers are serviced by one HSR line in a trip. As Figure 4 shows, the examples of intra-lines are 1-2-3-4-5-6-7, 1-11-12-13-14-8, and 7-8-9-10. Inter-line: passengers are serviced by two or more HSR

lines in a trip, passengers' origin and destination belong to one or more lines. For example, 5-6-7-8, 5-6-7-8-14, 7-8-14-13-12-11-1, and so on (Figure 4).

There are vague transportation organization modes in HSR in China so far. The average path length of the Chinese high-speed railway network is 1,403.36 km (17.89 h) according to actual lines. Thicker lines mean higher frequency of trains (Figure 5). Only a few lines with shorter distance and large passenger volume have high-frequency trains and approximate periodic timetables. Most lines in the network have lower frequency trains. There is a large percentage of inter-line trains in the HSR network.

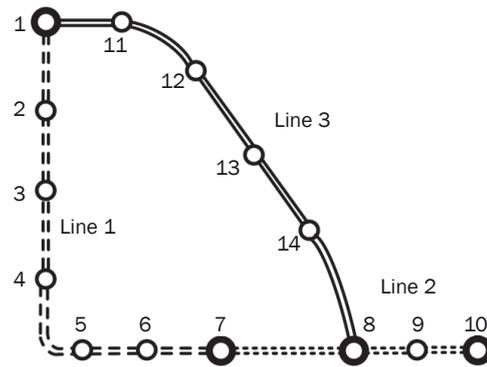


Figure 4 – Examples of intra-lines and inter-lines

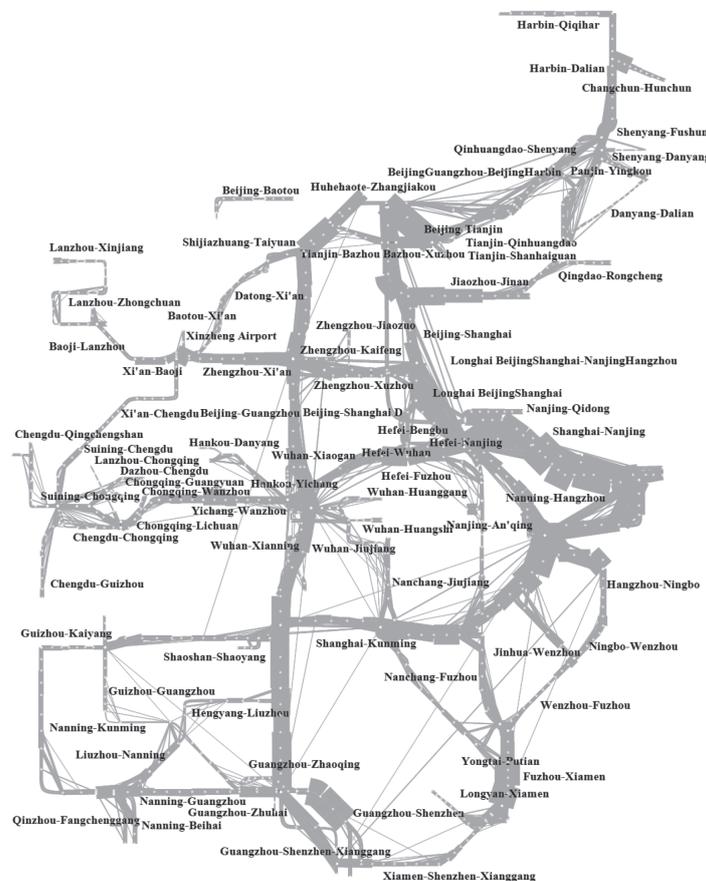


Figure 5 – Train service network of Chinese high-speed railway (2016)

1.3 Relevant research

Previous studies introduced some relevant methods detailing the transportation organization modes of HSR networks. They are classified into three categories: (1) transportation organization mode for specific high-speed lines, (2) passenger flow routing programming of high-speed railway network, and (3) passenger routing in line planning and timetabling.

Studies on transportation organization mode for specific high-speed lines. Many researchers explored the transportation organization modes of some typical high-speed lines, such as Shijiazhuang-Taiyuan HSR [1], Beijing-Shanghai HSR [2], Beijing-Tianjin HSR [3]. The inter-line mode, periodic or non-periodic, was discussed based on operation experience. There were three levels of rail services in the Netherlands' HSR [4]: IC, IR, and AR, which provided varied and fixed services based on a classified station system. A lot of features and experiences regarding the transportation organization modes of European [5] and Japanese [6-7] high-speed trains have been reported, including designing of train stops, train connections and transfers, combination modes of through trains and transfer trains via timetable mining. Train operation models were constructed to solve the relationship of transfer mode and nonstop mode [8-9] while determining the proportion of different train classes.

Passenger flow routing in HSR network. In previous studies, the object of passenger flow routing model would be minimizing train traveling time or maximizing the number of direct passengers. Enormous research efforts go into selecting the shortest path. Dijkstra's algorithm was used with a physical traffic network [9-13]. A model was put forward to pair nodes using an undirected graph, and undirected passenger routing was solved by a linear model which maximizes the number of direct passengers [14]. Passenger routing was transferred into a combinatorial optimization problem. The cross-entropy method was used to solve this problem [15]. Some academic achievements focused on passenger route choice behaviors and characteristics both in hubs and network. Surveys about quality service, adequate informing, quality of vehicles, line routes, and timetables were carried out [16-19].

Passenger routing in line planning and timetabling. Line planning and timetabling depends on passenger flow routing. The typical method of line planning is based on the line pool. Lines in the pool are selected according to the shortest routes. The shortest routes are selected as passenger flow routes based on maximum direct passengers. Claessens presented a mathematical programming model to solve the allocation of passenger lines [20]. This problem was based on maximizing the number of direct travelers. Schöbel

focused more on line planning. He found a set of paths for operating lines while aiming at minimizing transfers [21-22]. He used the Dantzig-Wolfe decomposition algorithm to solve the model. The sub-problem could be solved on the basis of shortest routes. Huiling Fu applied the line planning method from Schöbel to the Chinese HSR network to design a line plan considering cost-oriented and customer-oriented objectives [23-24].

In this work, we present a new approach for determining the transportation organization mode of an HSR network. We investigate the transportation organization mode via passenger flow routing in high-speed railway networks with respect to minimum total operation costs. The problem is solved through two-stages. First, given the railway infrastructure with the accompanying stations and passenger volumes between main OD pairs, the model determined operating cost-minimizing passenger flow routes. Then, the optimal results combining with route distances and line features are discussed to define the transportation organization mode of different transportation corridors.

2. BASIC HYPOTHESES

These hypotheses are made to depict passenger routing problems more logically and closer to reality. Cyclic annular railway can be avoided according to hypothesis 1. To describe the difference of passenger routes in an operation period more logically, hypothesis 2 was constructed. Hypothesis 3 is put forward to explain the determination of passenger flow for passenger routing.

1) The physical network is established based on cities, choosing one station per city. However, two or more stations in one city which cannot connect with each other via railway are included. Cyclic annular railway is avoided in this way.

Beijing South (L1) and Beijing (L3) belong to the city of Beijing, and two independent nodes are set up (Figure 6). Tianjin has three stations which can connect by lines in the Tianjin hub. Only one node is set up in this city.

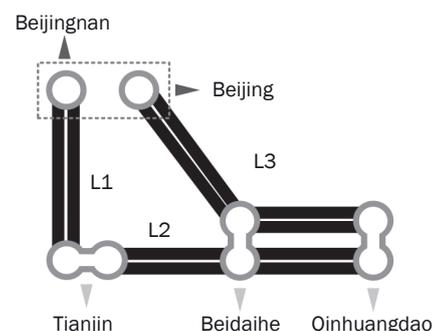


Figure 6 – An example of establishing a physical network

2) The construction cost and rolling stock cost are not considered because the costs are apportioned over a long operation period.

3) Passenger demand is divided into several pairs of nodes for cities with two or more stations, because the demand data of all pairs of nodes indicates the demand of two cities. As Figure 6 shows, the demand between Beijing and Qinhuangdao is averaged between Beijingnan-Qinhuangdao and Beijing-Qinhuangdao.

3. RESEARCH RESULTS AND PRESENTATION OF METHODS

A physical network is a finite, directed graph $PN=(N,E)$ with a node set N representing cities, and an edge set E which contains driving edges and transfer edges. For each node n . PN is constructed based on L-space; it means that for each edge $\{u,v\}$ there is a direct line from node u to station v . For each edge $\{u,v\}$ we assume the driving time t_{uv} , and the travel time differences between different lines t_{uv}^v are known. The

edge cost c_{uv} , is represented in three parts: driving cost, cost of travel time differences between different lines, and inter-line penalties per train.

Moreover, let $R \subseteq N \times N$ denote the set of all origin-destination pairs when q_{ij} means the sum of direct passengers. OD pairs can be served through the shortest route from i to j . The shortest route is calculated by Dijkstra's algorithm.

The definitions of main sets, parameters, and variables which are needed in the model are shown below (Table 1).

3.1 Enterprise model process

Two models (Table 2) are introduced to solve the passenger flow routing problem. M1 is established to solve maximum direct travelers based on the undirected graph, which only considers the number of connected lines for each node. The number of direct travelers is used in M2. The objective of M2 is minimum operation cost, while considering nodes capacity, edges capacity, and direct travelers. The solutions for paired nodes, direct routes, and total cost are found.

Table 1 – The definition of sets, parameters, and variables

Sets	R : set of OD pairs, all origin-destination pairs are contained
	V : set of demand of all origin-destination pairs
	K_i, K_j : the number of nodes i and j connected with
	$N_{R_{ij}}$: set of nodes of the route R_{ij}
	K_i^{up}, K_i^{dn} : the number of nodes i connected with upside and downside
Parameters	K_j^{up}, K_j^{dn} : the number of nodes j connected with upside and downside
	B_e, B_n : the bound of capacity for each node and edge
	q_{ij} : direct passengers that can be served from i to j through shortest routes
	c_{ij} : the operating cost from i to j through shortest routes per train
	y_{ij}^e : if the shortest route from i to j contains edge e or not
	z_{ij}^n : if the shortest route from i to j contains node n or not
	b_{ij}^e : if the route is chosen for ij pair & the route contains edge e or not
	w_{ij}^n : if the route is chosen for ij pair & the route contains node n or not
	L : allowed number of trains, 490 people per train
	θ : the proportion of direct passenger
M : the most number of direct passenger (solved by M1)	
Variables	x_{ij} : if i and j are paired or not, $x_{ij} \in \{0,1\}$

Table 2 – The distinction of M1 and M2

Model	M1	M2
Network	Undirected	Directed
Objective	Max. direct travelers	Min. operation cost
Constraints	a. The number of connected lines	a. The number of connected lines b. Nodes capacity c. Edges capacity d. Direct travelers
Results	a. Paired nodes b. Undirected routes c. The number of direct travelers	a. Paired nodes b. Directed routes c. Total cost

3.2 M1 for maximum direct travelers

To solve the passenger flow routing problem, the basic model was constructed [13]. Passenger flow routing with maximum direct travelers is presented. The M1 will pair nodes and select routes for maximum direct travelers based on the undirected graph.

$$\text{M1 objective } \max Q = \sum_{i \in PN} \sum_{j \in PN} q_{ij} \cdot x_{ij}$$

$$\text{s.t. } \sum_{j \in PN, i \neq j} x_{ij} \leq K_i, \quad \forall i, j \in PN \tag{1}$$

$$\sum_{i \in PN, i \neq j} x_{ij} \leq K_j, \quad \forall i, j \in NP \tag{2}$$

$$X_{ij} \begin{cases} 1 & \text{if } i \text{ and } j \text{ are paired} \\ 0 & \text{if } i \text{ and } j \text{ are not paired} \end{cases} \quad \forall i, j \in PN \tag{3}$$

$$q_{ij} > 0 \quad \forall i, j \in PN \tag{4}$$

Constraints 1 and 2 make sure that the number of routes through any node is not above the number of lines which the node connects. The variables $x_{ij} \in \{0,1\}$ in Constraint 3 are set to be 1 if and only if i and j are paired. Then Constraint 4 takes the passenger volume into account, which is greater than 0. The objective function we use is customer-oriented: we sum up the passenger volume $\sum_{i \in PN} \sum_{j \in PN} q_{ij} \cdot x_{ij}$ of a shortest route from i to j , we maximize direct travelers. It should be noted that q_{ij} means all demand can be transported through shortest route from i to j .

The routes with more nodes and longer traveling distance will be selected if maximum direct travelers with efficient and rapid transportation cannot be guaranteed at the same time. This solution is not suitable for a large-scale HSR network. The basic model disregards restriction of carrying capacity in railway sections, which will lead to capacity imbalance and complex transportation organization.

3.3 M2 for minimum total operation cost

We propose the passenger flow routing model with minimum operating cost (M2) while ensuring the proportion of direct passengers and considering capacity constraints. We established some special terms for the following research. (1) Inter-line trains. Trains run on two or more lines with different operation speeds. (2) Travel time differences between different lines. Trains run on lines with different maximum operating speeds due to inter-line operations. The time difference indicates the difference between trains running on higher speed lines and those running on lower speed lines.

The travel time difference due to operating speed is used in a route which contains edges with different operating speeds.

For the route A→B→C, the travel time differences between different lines will occur for edge BC. Travel time differences between different lines can be calculated by the formula $t_{BC}^{VACANT} = t_{BC}^{200} - t_{BC}^{250}$ (Figure 7), meaning that a train can run at 250 km/h regardless of the interline effect. Travel time differences between different lines of edge BC are between trains operating on 250 km/h and 200 km/h lines.

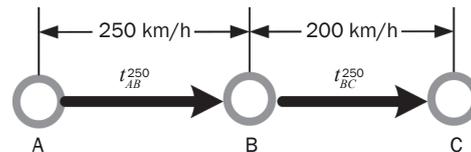


Figure 7 – An example of travel time differences between different lines

Constraints 1 and 2 in M1 will bring the new problem in a directed graph because constraints only restrict node-in and node-out. Passenger flow routes with the same direction on the same side for one node will be selected. Accordingly, the structure of the network needs to be amended. Four nodes are added into the nodes list for each station based on 2 directions of lines and 2 directions of nodes, node-in-up, node-in-down, nodes-out-up, and nodes-out-down, respectively.

A is the basic node form, B is the form of node-in and node-out, then the form changes into a four-part node in M2 (C) as shown in Figure 8.

Constraints 5 and 6 are proposed to make sure that the number of train routes through any node belonging to the network is not higher than that of the lines which the node connects.

$$\text{s.t. } \sum_{j \in PN, i \neq j} x_{ij} \leq K_i^{up}, \quad \sum_{j \in PN, i \neq j} x_{ij} \leq K_i^{dn} \tag{5}$$

$$\sum_{i \in PN} x_{ij} \leq K_j^{up}, \quad \sum_{i \in PN, i \neq j} x_{ij} \leq K_j^{dn} \tag{6}$$

Besides, Constraint 7 is put forward to guarantee a percentage of direct passengers because of the combination of direct and transfer transportation. The variable M is the maximum direct passengers, which results from M1. θ is the proportion of direct passengers and the range of θ is 0 to 1.

$$\sum_{i \in PN} \sum_{j \in PN} x_{ij} \cdot q_{ij} \geq \theta \cdot M \tag{7}$$

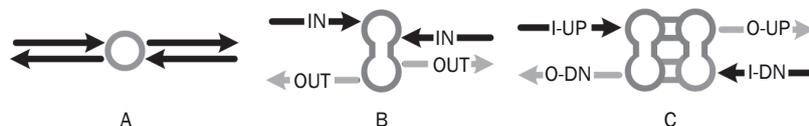


Figure 8 – The illustration of node expansion in PN

The operating cost per train contains route cost c_{ij}^t , cost of travel time differences between different lines c_{ij}^v , and penalty of inter-line c_{ij}^p (Equation 8). The route cost can be obtained by driving time multiplied by time cost because time cost is different at different operating speeds (Equation 9). For each route, the cost of travel time differences between different lines can be calculated by cost coefficient multiplied by travel time differences between different lines (Equation 10). Last, inter-line penalty is closely related to inter-line times (Equation 11).

$$c_{ij} = c_{ij}^t + c_{ij}^v + c_{ij}^p \tag{8}$$

$$c_{ij}^t = t_{ij}^{200} \cdot c^{200} + t_{ij}^{250} \cdot c^{250} \tag{9}$$

$$c_{ij}^v = t_{ij}^v \cdot c^v \tag{10}$$

$$c_{ij}^p = f_{ij} \cdot c^p \tag{11}$$

The cost of different train types can be obtained based on the CRH EMU technology parameters (Table 3).

Table 3 – Main cost parameters values [25]

Parameter	Values	Remarks
c^{200}	341 €/hour	Operating speed: 200 km/h
c^{250}	595 €/hour	Operating speed: 250 km/h and above
c^v	$2.8 \cdot 10^{-4}$ €/hour	Compare with maximum speed

$$\begin{aligned} \text{M2 objective: } & \min \sum_{i \in PN} \sum_{j \in PN} x_{ij} \cdot \frac{q_{ij}}{L} \cdot c_{ij} \\ \text{s.t. } & \sum_{i \in PN} \sum_{j \in PN} x_{ij} \cdot y_{ij}^e \cdot q_{ij} \leq B_e \end{aligned} \tag{12}$$

$$\sum_{i \in PN} \sum_{j \in PN} x_{ij} \cdot z_{ij}^n \cdot q_{ij} \leq B_n \tag{13}$$

$$q_{ij} \geq 0, B_e, B_n \geq 0 \tag{14}$$

$$\theta \in (0, 1] \tag{15}$$

$$X_{ij} \begin{cases} 1 & \text{if } i \text{ and } j \text{ are paired} \\ 0 & \text{if } i \text{ and } j \text{ are not paired} \end{cases} \quad \forall i, j \in PN \tag{16}$$

$$y_{ij}^e \begin{cases} 1 & \text{if the shortest route of } i \text{ to } j \\ & \text{contains edge } e \\ 0 & \text{if the shortest route of } i \\ & \text{to } j \text{ doesn't contain edge } e \end{cases} \quad \forall i, j, e \in PN \tag{17}$$

$$z_{ij}^n \begin{cases} 1 & \text{if the shortest route of } i \text{ to } j \\ & \text{contains node } n \\ 0 & \text{if the shortest route of } i \\ & \text{to } j \text{ doesn't contain node } n \end{cases} \quad \forall i, j, n \in PN \tag{18}$$

The boundaries of carrying capacity of edges and nodes are given in Constraints 12 and 13, respectively. Constraint 14 guarantees the passenger volume and

capacity are greater than 0. The variables $x_{ij} \in \{0, 1\}$ in Constraint 16 are set to 1 if and only if i and j are paired, and the variables $z_{ij}^e, w_{ij}^n \in \{0, 1\}$ in Constraints 17 and 18 are set to 1 if and only if edge e or node n is included in the shortest route from i to j . We sum up the operating cost of a shortest path from i to j , calculating the operating cost based on the number of operating trains, and we minimize the operating cost.

The introduced variables z_{ij}^e, w_{ij}^n will increase the complexity of the algorithm due to capacity constraints. The capacity judgment is updated in each iteration according to passenger assignment. The calculation process will be complicated and calculation time will be longer.

Other extensions variables are introduced to reduce the complexity of the algorithm. Constraints 12 and 13 translate into Constraints 19 and 20 by introducing new variables b_{ij}^e, w_{ij}^n . The variables b_{ij}^e, w_{ij}^n in Constraints 21 and 22 are set to 1 if i and j are paired, and the shortest route of i to j contains edge e or contains node n .

$$\sum_{i \in N} \sum_{j \in N} b_{ij}^e \cdot q_{ij} \leq B_e \tag{19}$$

$$\sum_{i \in N} \sum_{j \in N} w_{ij}^n \cdot q_{ij} \leq B_n \tag{20}$$

$$b_{ij}^e \begin{cases} 0 & \text{if the route isn't chosen for } i, j \text{ pair/} \\ & \text{the route doesn't contain edge } e \\ 1 & \text{if the route is chosen for } i, j \text{ pair \&} \\ & \text{the route contains edge } e \end{cases} \tag{21}$$

$$w_{ij}^n \begin{cases} 0 & \text{if the route isn't chosen for } i, j \text{ pair/} \\ & \text{the route doesn't contain node } n \\ 1 & \text{if the route is chosen for } i, j \text{ pair \&} \\ & \text{the route contains node } n \end{cases} \tag{22}$$

3.4 Algorithm processes

The algorithm processes are improved when new binary variables are applied. Passenger volume is assigned for the model after updating shortest routes.

A. Maximum direct volume for route R_{ij} from i to j is calculated with the formula $V_{R_{ij}} = \sum_{i \neq j, j=i+1}^j q_{ij}$. R_{ij} contains several nodes $N_{R_{ij}}$, and the passenger volume can be serviced for any two of these nodes paired in the direction of the route. The maximum of $\frac{N_{R_{ij}}(N_{R_{ij}} - 1)}{2}$ passenger volumes of OD pairs will be summed.

B. $V_{n-n+1} = \sum_{n \in N_{R_{ij}}, i \leq n < j} q_{nj}$ is used for passenger volume of each edge for each shortest route. $N_{R_{ij}} - 1$ edges will be assigned for route R_{ij} .

C. We calculate the passenger volume of each node for each shortest route by formula

$$V_{n-out-up} = \sum_{n \in N_{R_{ij}}, i \leq n < j} q_{nj}, V_{n-in-up} = \sum_{n \in N_{R_{ij}}, i < n \leq j} q_{in}, \text{ and } 2N_{R_{ij}} \text{ nodes will be assigned for route } R_{ij}.$$

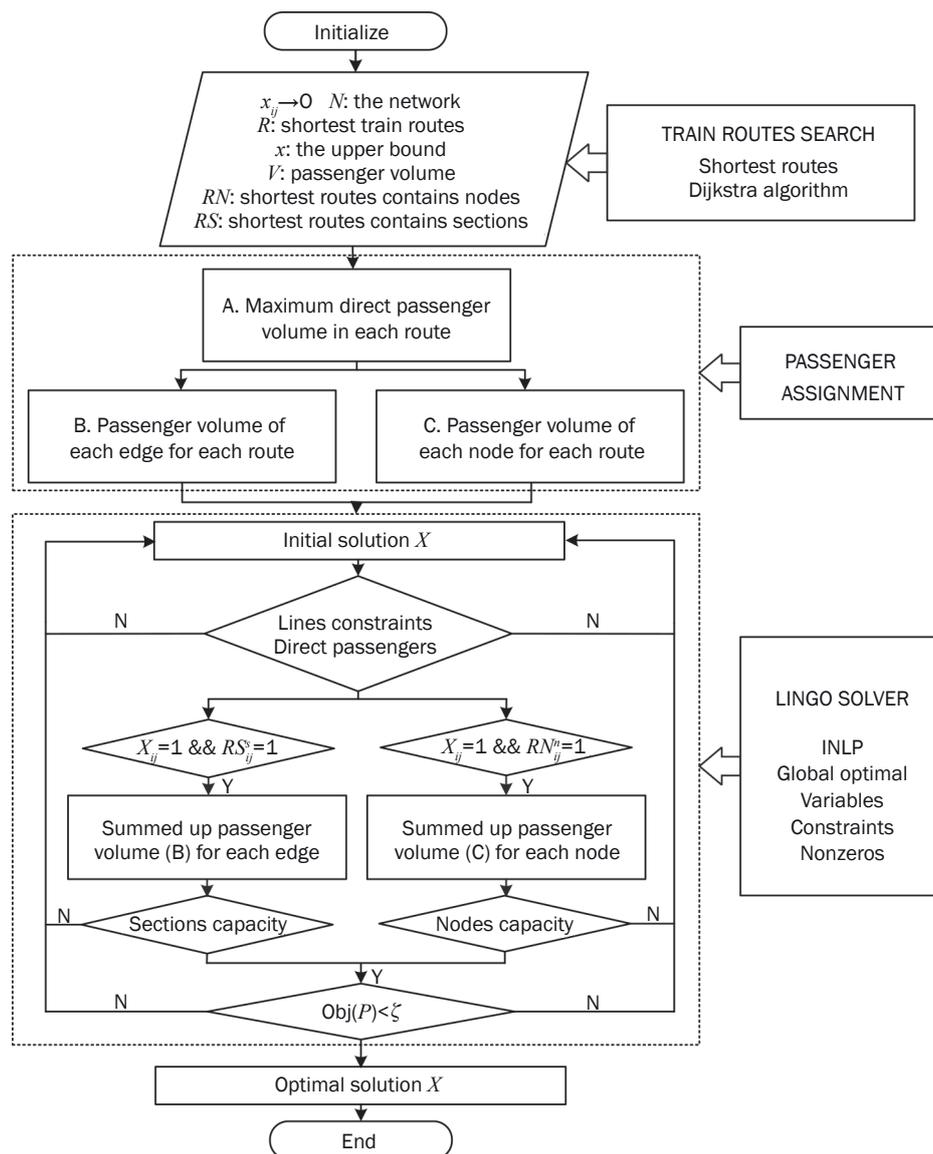


Figure 9 – The improved M2 algorithm process

The improved algorithm is shown in Figure 9. The initial solution X is calculated by the Lingo solver, while line constraints and direct passengers are judged, but the judgment of node and section capacity constraints should be solved based on passenger assignment. The passenger volume of each edge (B) and passenger volume of each node (C) for all selected routes are summed. If capacity constraints are not reached, a new solution will be applied, and the optimal solution will be found after many iterations.

A route with more passengers will be replaced by other routes if exceeding section capacity or nodes capacity limit in this iteration process. The process of iteration is the same as practical operation. The result shows a route which can provide a more direct service for OD pairs, and routes with less nodes are selected.

4. CASE STUDY

4.1 Case descriptions

Our approach is currently tested on a sub-network of the Northeast China railway network, which has seven high-speed railway lines (Table 4, Figure 10). The maximum operating speed is 300 km/h and the minimum operating speed is 200 km/h. The given PN contains 76 expanded nodes, 36 edges, and 341 origin-destination pairs. The demand data for the passenger flow routing problem is given by the origin/destination matrix, specifying the number of travelers per day between all pairs of nodes.

Then a model with 49,104 variables and 36,667 constraints is built. The shortest route algorithm and passenger assignment on each route are implemented

Table 4 – The basic information of PN

Numbers	Lines	Distance [km]	Speed [km/h]
1	Beijing-Tianjin	120	250
2	Tianjin- Qinhuangdao	350	250
3	Qinhuangdao-Shenyang	404	200
4	Beijing-Qinhuangdao	120	200
5	Harbin-Dalian	921	300
6	Changchun-Hunchun	456	200
7	Harbin- Qiqihar	282	250

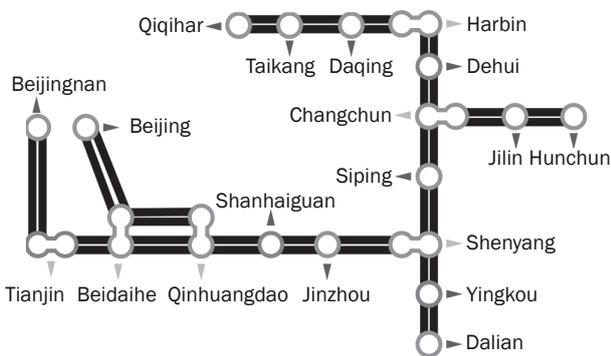


Figure 10 – The topological graph of the case

with C#, and Lingo 11.0 is used to solve the non-linear programming model with a global optimal algorithm. The solver performs on a PC with a 3.2 GHz, Intel™ Core™ i5-3470 processor and 16 GB of RAM running Windows 7 (64-bit).

4.2 Results and discussion

We discussed the model for different situations when parameters change.

The correlation between direct passengers and total cost

The direct passenger volume grows in inverse proportion to the increase of total cost. The more direct passenger demand is satisfied, the more cost it produces. The growth of total costs accelerates when the proportion of direct passengers is greater than 80%. A rapid increase of total cost occurs when parameter θ is greater than 85% (Figure 11).

The characteristics of routes

The different train routes are optimized. The characteristics of the routes are summarized when parameter θ changes. The average nodes number, distance, and inter-line times increase with the proportion of direct passengers (Table 5). The rate of increase in these statistical values slows down when θ is greater than 90% (Figure 12). More routes are selected when the

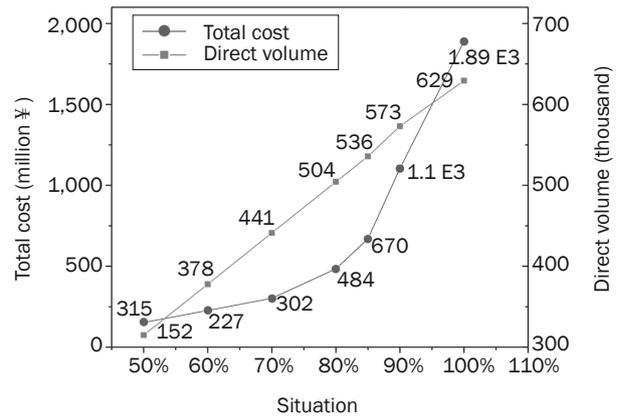


Figure 11 – The correlation between direct passenger and total cost

Table 5 – The statistical value of different situation

Situation	The number of routes	Avg. nodes	Avg. distance	Avg. inter-line times
100%	9	5.000	622.444	0.667
90%	9	4.778	570.111	0.444
85%	11	4.545	476.545	0.364
80%	15	3.333	348.200	0.267
70%	21	2.550	236.667	0.190
60%	22	2.409	203.273	0.182
50%	20	2.350	178.400	0.250

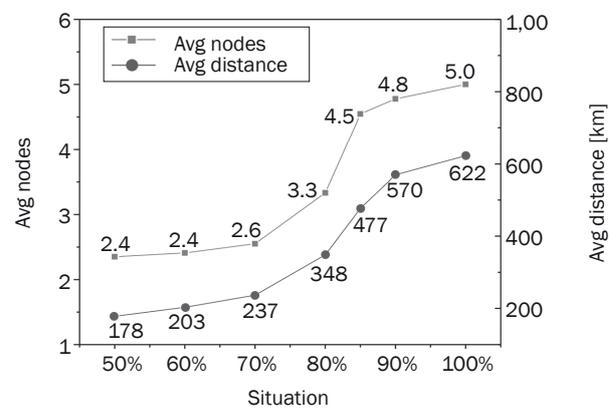


Figure 12 – The change of average nodes and average distance when parameter θ changed

proportion of direct passengers decreases because routes with more nodes can satisfy direct passenger volumes better.

The transportation organization mode

The longer distance routes are divided into several parts with the reducing proportion of direct passengers. The longest distance of routes is 1,565 km (Figure 13a). The total cost and difficulty of operation increase rapidly. The longest routes are from Beijing

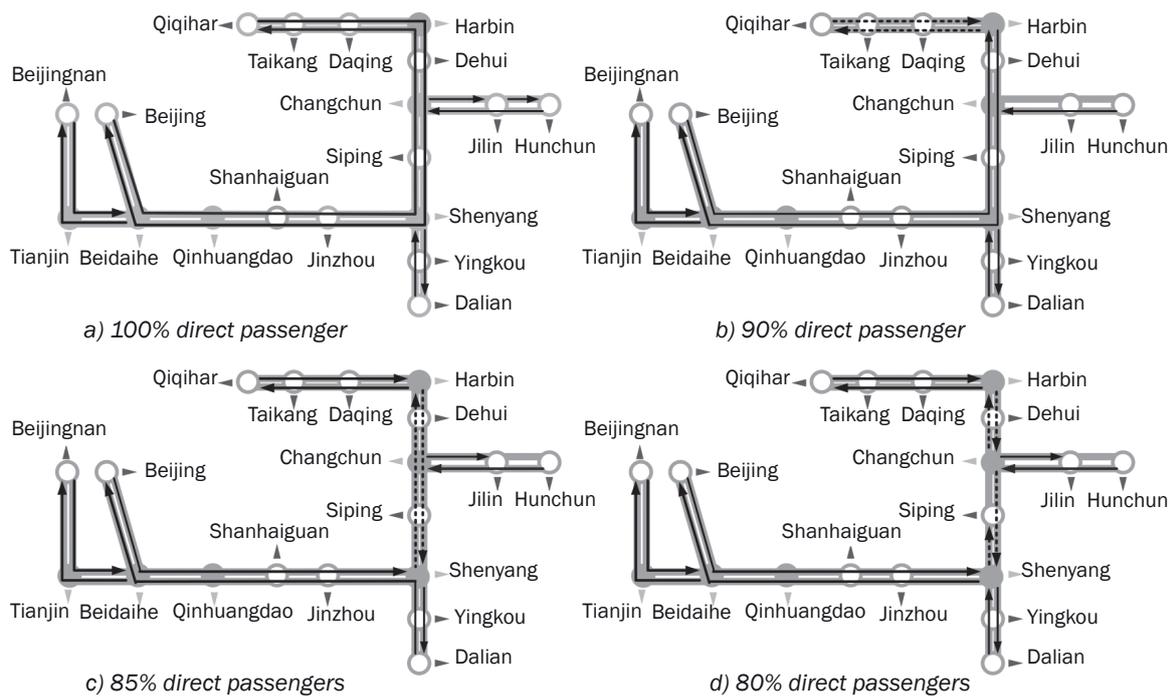


Figure 13 – The illustration of paired nodes and routes in different situation

to Harbin or Harbin to Beijing (Figure 13b). The main change in Figure 13c is that the longest routes in Figure 13b are divided into two parts. The route from Beijing to Dalian is the only longest route when parameter θ is 85% (Figure 13c). The optimal result shows an imbalance in the spatial distribution of passenger flow. The optimal result in Figure 13d is unrealistic because the influence of the Harbin-Dalian HSR is severely attenuated.

The transportation organization mode of the network in the case is demonstrated by the exiting line plan (Figure 14). It is similar to the optimal result of 90% direct passengers, but the total cost is 12.5% higher in comparison.

The optimal result using the second optimal model (M2) is confirmed as the result of making the transportation organization mode. The transportation organization modes of transport corridors are shown in Table 6. Intra-line or inter-line and periodic or non-periodic modes are determined based on route distances and line features. Inter-line trains will operate in corridors if routes are composed of several lines. Non-periodic mode will be adopted in corridors if inter-line routes have longer distances.

Passengers can finish trips only by taking one train when both origin and destination are in a result route. Passengers whose origin and destination is situated in different routes need a transfer to finish their trips based on results (Table 6).

Table 6 – The operation mode of main corridors in sub-network

OD Pairs	Distance [km]	Intra-line / Inter-line	Periodic/Non-periodic
Changchun-Jilin	111	Intra-line	Periodic
Hunchun-Changchun	471	Intra-line	Periodic
Dalian-Shenyang	400	Intra-line	Non-periodic
Beijing-South-Beidaihe	365	Inter-line	Periodic
Beidaihe-Beijing-South	365	Inter-line	Periodic
Qiqihar-Harbin	312	Intra-line	Periodic
Harbin-Qiqihar	312	Intra-line	Periodic
Shenyang-Harbin	550	Intra-line	Periodic/Non-periodic
Harbin-Shenyang	550	Intra-line	Periodic/Non-periodic
Shenyang-Beijing	703	Inter-line	Non-periodic
Beijing-Dalian	1103	Inter-line	Non-periodic

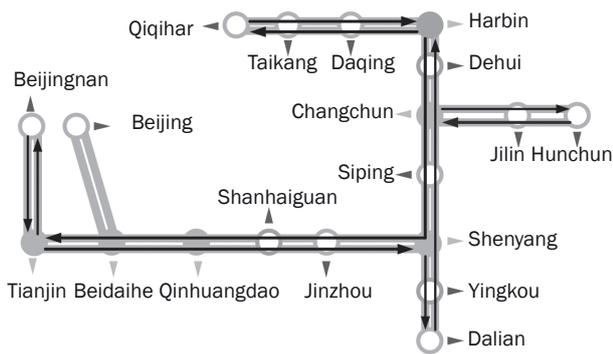


Figure 14 – The illustration of paired nodes and routes of the actual network

5. CONCLUSIONS

We developed an integer non-linear programming model for solving the passenger flow routing problem that minimizes the operating costs, including traveling cost, cost of travel time differences between different lines, and penalty for the inter-line. Two binary variables were introduced to simplify the flow diagrams algorithm. We discussed optimal results in different situations, and the transportation organization modes of a sub-network of the Northeast China railway network were determined.

Different transportation organization modes in network and unbalanced passenger flow are described in the passenger routing model. The method can be used to organize passenger flow in a large-scale HSR network at different development stages. We can determine the transportation organization mode of a region according to passenger routes.

On the one hand, the results provide a line pool reference which adapts to the complex operation environment for HSR network line planning. Network lines will be more effective considering different transportation organization modes, especially for multiple line speeds, unbalanced passenger flows, and different grades of stations.

On the other hand, research on passenger organization in hubs can be improved based on passenger routing models. Passenger organization in a hub containing multiple stations is also complex. A passenger routing model provides new ideas for the division of labor at the station within a hub.

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基于不同运输模式的高速铁路网络客流路径优化模型

摘要

为了适应大规模的高速铁路网络特征, 根据不同运输组织模式进行合理的客流路径规划是提高高速铁路网络化运营效率的重要措施。论文以旅客运输模式为基础, 提出客流路径模型优化高速铁路网络的客流路径。模型中以运营费用最小为研究目标, 运营费用包含出行成本、由于线路运营速度不同产生的差异化运行时间费用和跨线惩罚费用。为了更好地适应有向图的场景对网络进行重构, 利用四个节点表示一个车站, 四个节点包括进站上行、进站下行、出站上行和出站下行。以该网络为基础提出整数非线性模型用以解决客流路径优化问题, 模型通过引入二进制变量优化模型计算流程。构建了中国东北地区的高速铁路网为案例对模型进行验证, 通过Lingo11.0可得到最优解, 根据得到客流量、运输距离和基础设施条件对网络的运输组织模式进行探讨。

关键词

高速铁路; 运输计划; 运营模式; 客流路径; 整数非线性规划模型

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